

# **International Economics B**

## **2. Basics in noncooperative game theory**

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## What is “game theory”?

- Consider situations in which more than one person (“agents”) interact with each other.
  - One agent’s decision affects other agents’ well-being (“payoff”).
- In the presence of such **strategic interactions**, study
  - How does each agent make a decision?
  - What outcome is produced as a consequence of the agents’ behavior?

## Basic concepts in noncooperative game theory

- **Noncooperative game theory:** Players (=agents who participate in a game) make decisions **independently**.
  - cf. Cooperative game theory: Assuming the possibility of external enforcement of cooperative behavior (e.g., through contract law)
- **Representations of games**
  - Normal-form (or Strategic-form) game: Useful to describe games in which players **simultaneously** make decisions.
  - Extensive-form game: Useful to describe games in which **sequentially** make decisions.

## Assumptions:

- All players have the accurate knowledge about the structure, rules, and payoffs of the game.
- Perfect information: Each player has all the information concerning the actions taken by other players earlier in the game that affect the player's decision about which action to choose at a particular time.
  - cf. Imperfect information

## Formal definition of a normal-form game

### Definition

A **normal form game** is described by:

- 1 A set of **players**:  $I \equiv \{1, 2, \dots, N\}$ .
- 2 An **action set** of each player:  $a^i \in A^i$ ,  
 $A^i = \{a_1^i, a_2^i, \dots, a_{k_i}^i\}$ , which is the set of all actions available to player  $i$ ;

An **outcome** of the game:  $a \equiv (a^1, a^2, \dots, a^i, \dots, a^N)$ , which is a list of the actions chosen by each player.

- 3 A **payoff function** of each player:  
 $\pi^i(a) = \pi^i(a^1, \dots, a^N)$ .

## Examples of normal-form games

### Example 1: Peace-War game

		Country 2	
		WAR	PEACE
Country 1	WAR	1    1	3    0
	PEACE	0    3	2    2

In this game,

- Set of players: {Country 1, Country 2}
- Each player's action set:  $A^1 = A^2 = \{\text{WAR}, \text{PEACE}\}$
- Possible outcomes of the game: (W, W), (W, P), (P, W), (P, P)
- Payoff: when  $a = (\text{W}, \text{P})$ ,  $\pi^1(a) = 3$  and  $\pi^2(a) = 0$

## Solving the game

- An **equilibrium** of a game: How will the game end up from the all possible outcomes?
- The most commonly used solution concept is the Nash equilibrium.
  - Nash, J. (1951), "Non-Cooperative Games," *Annals of Mathematics* 54(2), 286-295.
- A set of strategies is a Nash equilibrium if no player can do better by unilaterally changing his/her strategy.

Denote the list of actions by players except player  $i$  by  $a^{-i}$ ,  
i.e.,

$$a^{-i} \equiv (a^1, \dots, a^{i-1}, a^{i+1}, \dots, a^N).$$

$\Rightarrow$  An outcome  $a$  can be expressed as the list of player  $i$ 's  
action and actions of players other than player  $i$ :  
 $a = (a^i, a^{-i})$ .



## Formal definition of a Nash equilibrium

### Definition

An outcome  $\hat{a} = (\hat{a}^1, \hat{a}^2, \dots, \hat{a}^N) \in A^1 \times A^2 \times \dots \times A^N$  is a **Nash equilibrium** if no player has an incentive to deviate from  $\hat{a}^i$  provided that all other players do not deviate from  $\hat{a}^{-i}$ .

Formally, for every player  $i = 1, 2, \dots, N$ ,

$$\pi^i(\hat{a}^i, \hat{a}^{-i}) \geq \pi^i(a^i, \hat{a}^{-i}) \quad \forall a^i \in A^i.$$

## Best-response function and Nash equilibrium

### Definition

The **best-response function** of player  $i$  is the function  $R^i(a^{-i})$  that assigns, for given actions  $a^{-i}$  of other players, an action  $a^i = R^i(a^{-i})$  that maximizes player  $i$ 's payoff  $\pi^i(a^i, a^{-i})$ .

### Theorem

If  $\hat{a} = (\hat{a}^1, \dots, \hat{a}^N)$  is a Nash equilibrium outcome, then  $\hat{a}^i = R^i(\hat{a}^{-i})$  holds for every player  $i$ .

## Example 1 (Peace-War game):

		Country 2	
		WAR	PEACE
Country 1	WAR	1    1	3    0
	PEACE	0    3	2    2

- Country 1's best-response function:

$$R^1(a^2) = \begin{cases} \text{WAR} & \text{if } a^2 = \text{WAR} \\ \text{WAR} & \text{if } a^2 = \text{PEACE} \end{cases}$$

- Country 2's best-response function:

$$R^2(a^1) = \begin{cases} \text{WAR} & \text{if } a^1 = \text{WAR} \\ \text{WAR} & \text{if } a^1 = \text{PEACE} \end{cases}$$

**(WAR, WAR) is a (unique) Nash equilibrium.**

**The procedure for finding a Nash equilibrium:**

- ① Calculate the best-response function of each player.**
- ② Find outcomes that lie on the best-response functions of all players.**

**Not all games have a unique Nash equilibrium.**

- Multiple Nash equilibria**
- Nonexistence of a Nash equilibrium**

## Multiple Nash equilibria

### Example 2: Battle of the sexes

		Rachel	
		OPERA ( $\omega$ )	FOOTBALL ( $\phi$ )
Jacob	OPERA ( $\omega$ )	2      1	0      0
	FOOTBALL ( $\phi$ )	0      0	1      2

- Both of them gain a higher utility if they go together to one of these events: “coordination game”

There are two Nash equilibria: (OPERA, OPERA) and (FOOTBALL, FOOTBALL).

## Nonexistence of a Nash equilibrium

### Example 3: Battle of the sexes after 30 years of marriage

		Rachel	
		OPERA ( $\omega$ )	FOOTBALL ( $\phi$ )
Jacob	OPERA ( $\omega$ )	2      0	0      2
	FOOTBALL ( $\phi$ )	0      1	1      0

- J wants to be with R, but R wants to be alone.

There is no (pure strategy) Nash equilibrium.

- A Nash equilibrium in mixed strategy exists.
  - Mixed strategy: an assignment of a probability to each pure strategy (i.e., players randomly chooses each pure strategy).

## Extensive-form games

- Games with dynamic interactions: represented by **extensive forms** (game trees).
- The extensive form of a game is a complete description of:
  - 1 The set of players;
  - 2 Who moves when and what their choices are;
  - 3 What players know when they move;
  - 4 The players' payoffs as a function of the choices that are made.

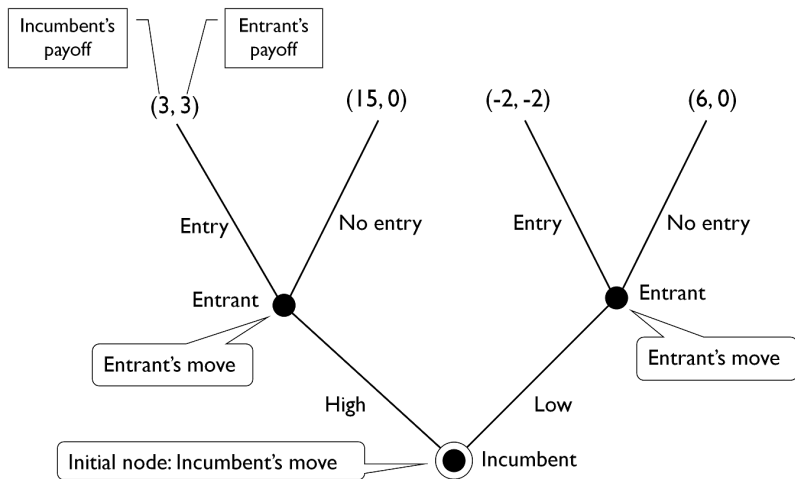
## Definition

- An **extensive form game** consists of:
  - 1 A **game tree** containing
    - a starting node, other decision nodes, terminal nodes, and
    - branches linking each decision node to successor nodes.
  - 2 A list of **players**  $i = 1, 2, \dots, N$ .
  - 3 For each decision node, the name of the player entitled to choose an action.
  - 4 Each player  $i$ 's **action set** at each decision node.
  - 5 Each player  $i$ 's **payoff** at each terminal node.



## Example 4: Entry deterrence

- **Players:** an incumbent firm & a new entrant
- **The order of play:**
  - ① Incumbent determines the price of its product: “High” or “Low”.
  - ② Entrant decides whether “Entry” or “No entry”.
- **Each player’s payoff:**
  - “High” & “Entry” → Both firms earn 3 million\$.
  - “High” & “No entry” → Incumbent = 15 million\$, Entrant = 0.
  - “Low” & “Entry” → Both firms lose money (-2 million\$).
  - “Low” & “No entry” → Incumbent = 6 million\$, Entrant = 0.



## Strategies and outcomes in extensive form games

### Definition

A **strategy** for player  $i$ ,  $s^i$ , is a complete plan (list) of actions, one action for each decision node that the player is entitled to choose an action.

- Not what the player does at a single specific node but is a list of what the player does at every node where the player is entitled to choose an action.

## In Example 4 (Entry deterrence),

- **Incumbent:** One decision node (= initial node)  
⇒ **Strategy:** H or L
- **Entrant:** Two decision nodes (left and right)  
⇒ **Specification of the precise action taking at each node.**
  - E at both nodes → **Strategy:** (E, E)
  - E at left & NE at right → (E, NE)
  - NE at left & E at right → (NE, E)
  - NE at both nodes → (NE, NE)
- **Possible outcomes:**  
(H, (E, E)), (H, (E, NE)), (H, (NE, E)), (H, (NE, NE)),  
(L, (E, E)), (L, (E, NE)), (L, (NE, E)), (L, (NE, NE))

## Solving the game

**Solution concept: Subgame perfect (Nash) equilibrium**

- **A refinement of the Nash equilibrium concept proposed by Selten (1965).**
  - Selten, R. (1965), “Spieltheoretische Behandlung eines Oligopolmodells mit Nachfrageträgheit,” *Zeitschrift für die gesamte Staatswissenschaft* 121, 301–24 and 667–89.
- **An equilibrium such that players’ strategies constitute a Nash equilibrium in every subgame of the original game.**
  - **Subgame:** A decision node from the original game along with the decision nodes and terminal nodes directly following this node.
- **Subgame perfect equilibria eliminate noncredible threats.**

- Rewrite the game in a normal form:

		Entrant							
		(E,E)		(E,NE)		(NE,E)		(NE,NE)	
Incumbent	H	3	3	3	3	15	0	15	0
	L	-2	-2	6	0	-2	-2	6	0

- Two Nash equilibria:  $(H, (E, E))$  and  $(L, (E, NE))$
- However,  $(H, (E, E))$  includes noncredible threats.
  - $(E, E)$ : The entrant chooses “E” regardless of the incumbent’s strategy.
  - Once the incumbent chooses “L”, the entrant will prefer “NE” to “E”.

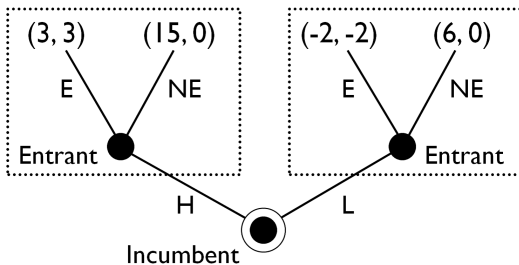
## Subgame

### Definition

- A **subgame** is a decision node from the original game along with the decision nodes directly following this node.
- A subgame is called a **proper subgame** if it differs from the original game.

**Example 4 (Entry deterrence game) has three subgames:**

- A subgame after the incumbent chooses H
- A subgame after the incumbent chooses L
- The original game





## Formal definition of the subgame perfect equilibrium

### Definition

An outcome is a **subgame perfect equilibrium** if it induces a Nash equilibrium in every subgame of the original game.

- A subgame perfect equilibrium outcome is also a Nash equilibrium of the original game.

## Backward induction and the SPE

**Method for finding the SPE outcome: Backward induction**

- 1 Find the NE of in the subgames leading to the terminal nodes (i.e., optimal strategy of the player who makes the last move of the game).**
- 2 Find the NE for the subgames leading to the subgames leading to the terminal nodes (i.e., optimal choice of the next-to-last moving player), taking as given the NE actions played in the last subgames.**
- 3 Continuing to solve in this way backwards in time until all players' actions have been determined.**

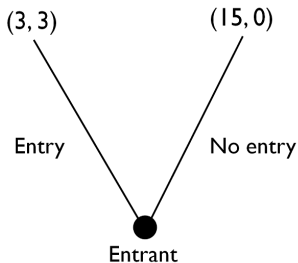
**Find the subgame perfect equilibrium in Example 4.**

**① 2nd stage (entrant's move): Two subgames**

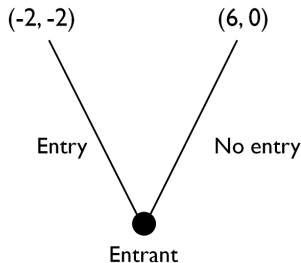
- A subgame after the incumbent chooses H
- A subgame after the incumbent chooses L

**⇒ In each subgame, the entrant chooses its optimal action: E or NE.**

**② 1st stage (incumbent's move): Taking the entrant's optimal choice in the 2nd stage, the incumbent chooses H or L.**

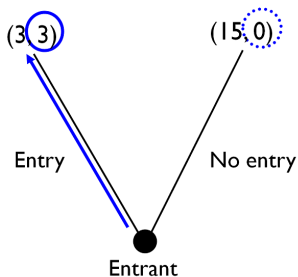


A subgame after "High"

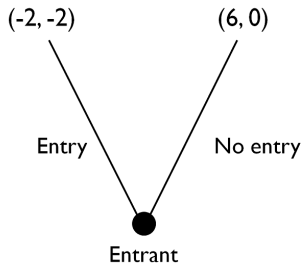


A subgame after "Low"

⇒ Entrant's optimal strategy:

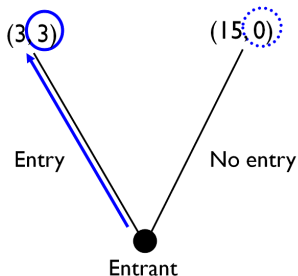


A subgame after “High”

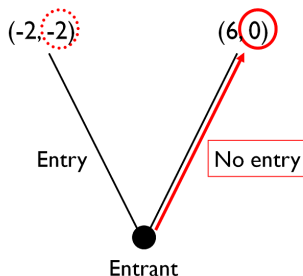


A subgame after “Low”

⇒ Entrant's optimal strategy:

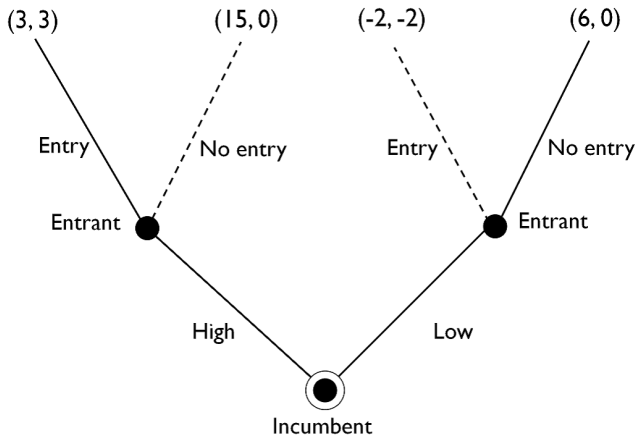


A subgame after "High"

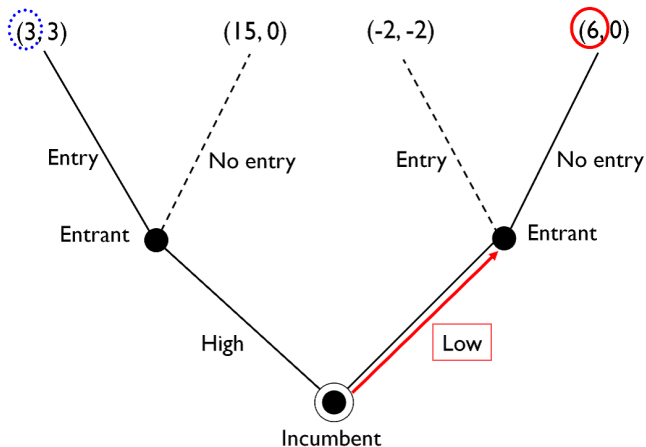


A subgame after "Low"

⇒ Entrant's optimal strategy: (E, NE)



⇒ Incumbent's optimal strategy:



⇒ Incumbent's optimal strategy: **L**



**In Example 4 (Entry-deterrence game),**

- **Subgame perfect equilibrium:  $(L, (E, NE))$**
- **Outcome of the game: The incumbent chooses “Low”, and then the entrant chooses “No entry”.**

## References

- **Mas-Colell, A., M.D. Whinston, and J.R. Green (1995),** *Microeconomic Theory*, **Oxford University Press.**
- **Shy, O. (1996),** *Industrial Organization: Theory and Applications*, **The MIT Press.**
- **Gibbons, R. (1992),** *Game Theory for Applied Economists*, **Princeton University Press.**